

1/2

FIG 1

Input: (implicit: topology, routing, budgets)

 $B_{hot} := B$ while $B_{hot} \neq \emptyset$ doselect $b^* \in B_{hot}$ with the largest blocking probability $c_U^{inc} := 1$ if $(\forall l \in \mathcal{E}: c_U^{free}(l) \geq c_U^{inc}(l, b^*))$ then $c_U[b^*] := c_U[b^*] + c_U^{inc}$

else

 $B_{hot} := B_{hot} \setminus b^*$

end if

end while

Output: assignment of portions of transmission capacity

 $c_U[b], b \in B$

FIG 2

Input: Link l (implicit: topology, routing, budgets)if $|\{b: b \in B_{hot} \wedge u(l, b) > 0\}| > 0$ then select $b^* \in B_{hot} : u(l, b^*) > 0$ with the largest blocking probability $c_U^* := \lfloor q(l) \cdot a(b^*) \rfloor$ $p_b^* := p_b(a(b^*), c_U[b^*] + c_U^*)$ for all $b \in \{b: b \in B_{hot} \wedge u(l, b) > 0\}$ do $c_U^b := \lfloor q(l) \cdot a(b) \rfloor$ $p_b^b := p_b(a(b), c_U[b] + c_U^b)$ while $p_b^* < p_b^b$ do $c_U^* := \lfloor q^{dec} \cdot c_U^* \rfloor$ $p_b^* := p_b(a(b^*), c_U[b^*] + c_U^*)$

end while

end for

else

 $c_U^* := 0$

end if

Output: link capacity increment c_U^*

2/2

FIG 3

Input: (implicit: topology, routing, budgets)

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    for all  $l \in \mathcal{E}$  do
         $c_U^{inc}[l] := \text{CapInc}(l)$ 
    end for
     $B_{hot} := B$ 
    while  $B_{hot} \neq \emptyset$  do
        select  $b^* \in B_{hot}$  with the largest blocking probability

         $c_U^{inc} := \max(1, \min_{l \in \mathcal{E}: u(l,b) > 0} c_U^{inc}[l])$ 
        if  $(\forall l \in \mathcal{E}: c_U^{free}(l) \geq c_U^{inc} \cdot u(l, b^*))$  then
             $c_U[b^*] := c_U[b^*] + c_U^{inc}$ 
        else
             $B_{hot} := B_{hot} \setminus b^*$ 
        end if
        for all  $l \in \mathcal{E}$  do
            if  $u(l, b^*) > 0$  then
                 $c_U^{inc}[l] := \text{CapInc}(l)$ 
            end if
        end for
    end while

```

Output: assignment of portions of transmission capacity

 $c_U[b], b \in B$

FIG. 3: Algorithm 3